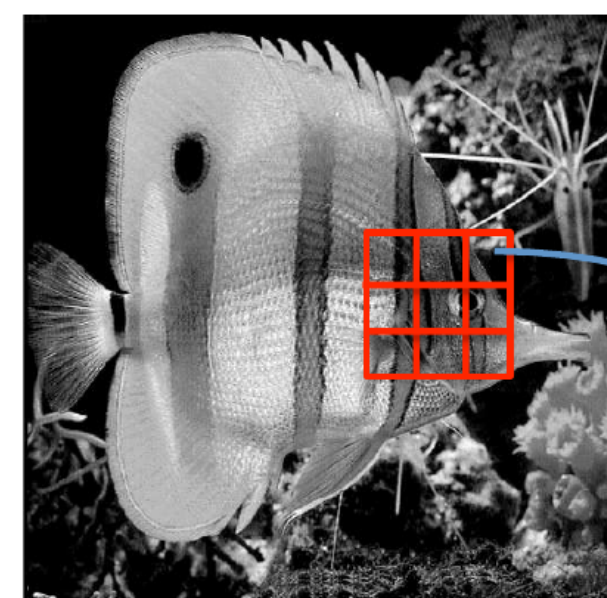
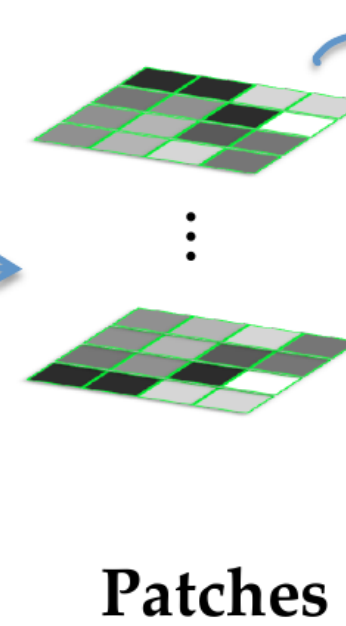


Motivation: Dictionary Learning

Given Y , find (A, X) such that $Y \approx AX$, with X as sparse as possible.

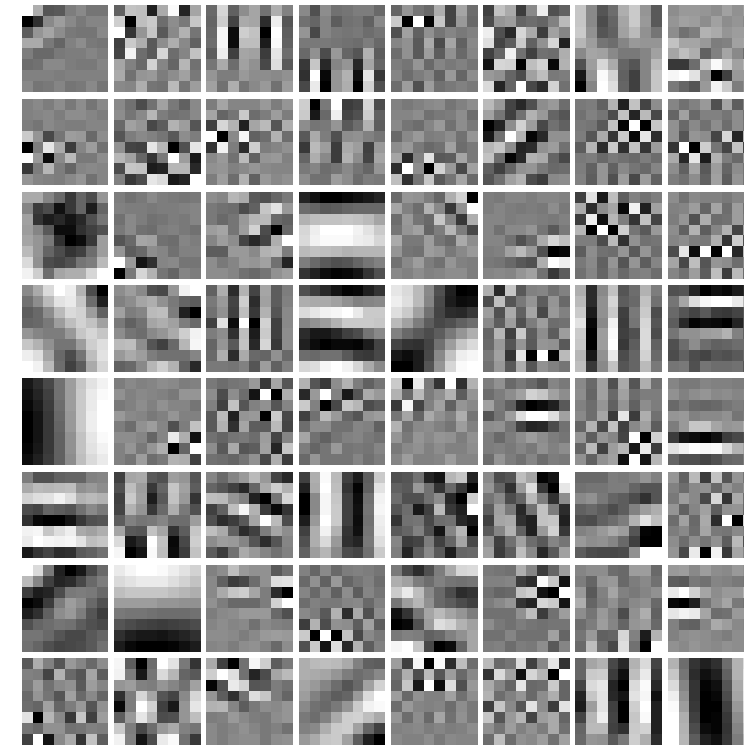


An image



Patches

$$Y \in \mathbb{R}^{n \times p}$$



Very successful in classical image processing, visual recognition, compressive signal requisition, and recently into deep architectures for signal classification ... **limited theoretical understanding.**

Problem: Dictionary Recovery

Given $n \times p$ data matrix $Y = A_0 X_0$ with X_0 sparse, recover A_0 and X_0 .

- ▶ Even A_0 is known, seeking X_0 is generally hard (sparse recovery problem!)
- ▶ Recovery only up to sign, permutation, and scale (as $A_0 X_0 = A_0 \Pi \Sigma * \Sigma^{-1} \Pi^* X_0$ for any permutation Π and full rank diagonal Σ)
- ▶ Due to the symmetry, **hard to convexify** the problem!

We focus on the case A_0 is complete (square and invertible), and the coefficients X_0 obeys a **Bernoulli-Gaussian model**: $[X_0]_{ij} = V_{ij} B_{ij}$, with $V_{ij} \sim \mathcal{N}(0, 1)$ and $B_{ij} \sim \text{Ber}(\theta)$.

Main Result (informal)

For any $\theta \in (0, 1/3)$, given $Y = A_0 X_0$ with A_0 a complete dictionary and $X_0 \sim_{i.i.d.} \text{BG}(\theta)$, there is a **polynomial-time algorithm** that recovers A_0 and X_0 with high probability (at least $1 - O(p^{-6})$) whenever

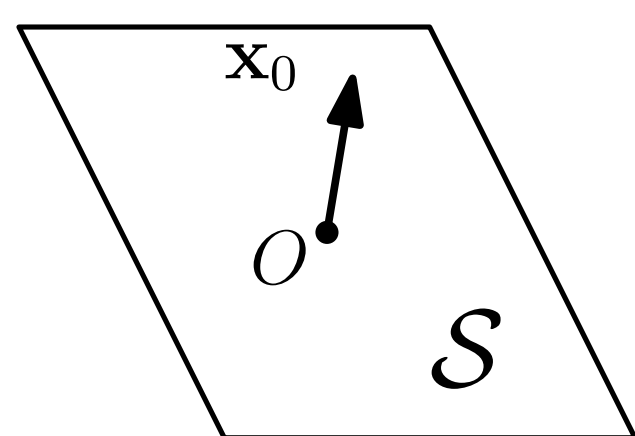
$$p \geq p_*(n, 1/\theta, \kappa(A_0), 1/\mu)$$

for a fixed polynomial $p_*(\cdot)$, where $\kappa(A_0)$ is the condition number of A_0 and μ is a smoothing parameter which can be set as $\mu = cn^{-5/4}$.

Main Ingredients I - A Nonconvex Formulation

- ▶ When A_0 is complete, $\text{row}(Y) = \text{row}(X_0)$.
- ▶ Rows of X_0 are sparse vectors in $\text{row}(Y)$. When $p \geq \Omega(n \log n)$, they are also the sparsest ones! [Spielman et al '12]

Find sparsest vectors in a given linear subspace ...



Natural formulation:

$$\text{minimize } \|q^* Y\|_0 \quad \text{subject to } q \neq 0.$$

Convex relaxation: [Spielman et al'12]

$$\text{minimize } \|q^* Y\|_1 \quad \text{subject to } \|q^* Y\|_\infty = 1.$$

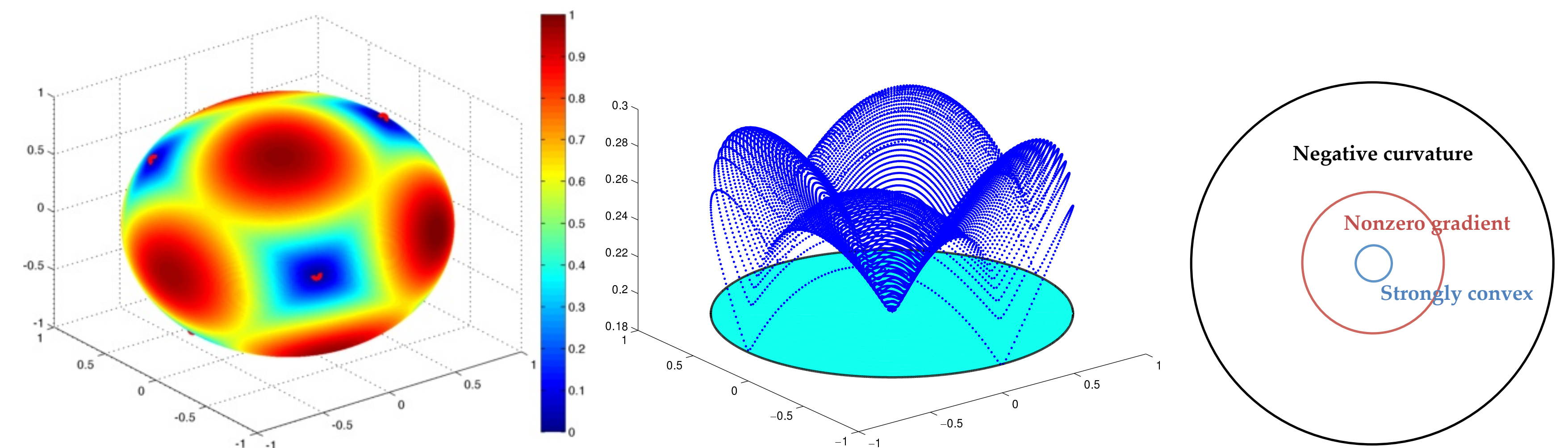
Convex relaxation is known to break down when each column of X_0 contains more than $O(\sqrt{n})$ nonzeros. We look at a nonconvex "relaxation":

$$\text{minimize } f(q; Y) \doteq \frac{1}{p} \sum_{k=1}^p h_\mu(q^* y_k) \quad \text{subject to } \|q\|_2 = 1.$$

with $h_\mu(\cdot)$ a smooth approximation to the $|\cdot|$ function:

$$h_\mu(z) = \mu \log \left(\frac{\exp(z/\mu) + \exp(-z/\mu)}{2} \right) = \mu \log \cosh(z/\mu)$$

Main Ingredients II - A Glimpse into High-dimensional Geometry



Theorem: (informal) Suppose $A_0 = I$ and hence $Y = A_0 X_0 = X_0$. For $\theta \in (0, 1/2)$ and μ sufficiently small $O(n^{-c})$, whenever $p \geq \frac{C}{\mu^2 \theta^2} n^3 \log \frac{n}{\mu \theta}$, the following hold simultaneously w.h.p.:

$$\begin{aligned} \nabla^2 g(w; X_0) &\succeq \frac{c_* \theta}{\mu} I & \forall w \quad \text{s.t.} \quad \|w\| &\leq \frac{\mu}{4\sqrt{2}}, \\ \frac{w^* \nabla g(w; X_0)}{\|w\|} &\geq c_* \theta & \forall w \quad \text{s.t.} \quad \frac{\mu}{4\sqrt{2}} &\leq \|w\| \leq \frac{1}{20\sqrt{5}}, \\ \frac{w^* \nabla^2 g(w; X_0) w}{\|w\|^2} &\leq -c_* \theta & \forall w \quad \text{s.t.} \quad \frac{1}{20\sqrt{5}} &\leq \|w\| \leq \sqrt{\frac{4n-1}{4n}}, \end{aligned}$$

and the function $g(w; X_0)$ has **exactly one local minimizer** w_* over the open set $\Gamma \doteq \left\{ w : \|w\| < \sqrt{\frac{4n-1}{4n}} \right\}$, which satisfies

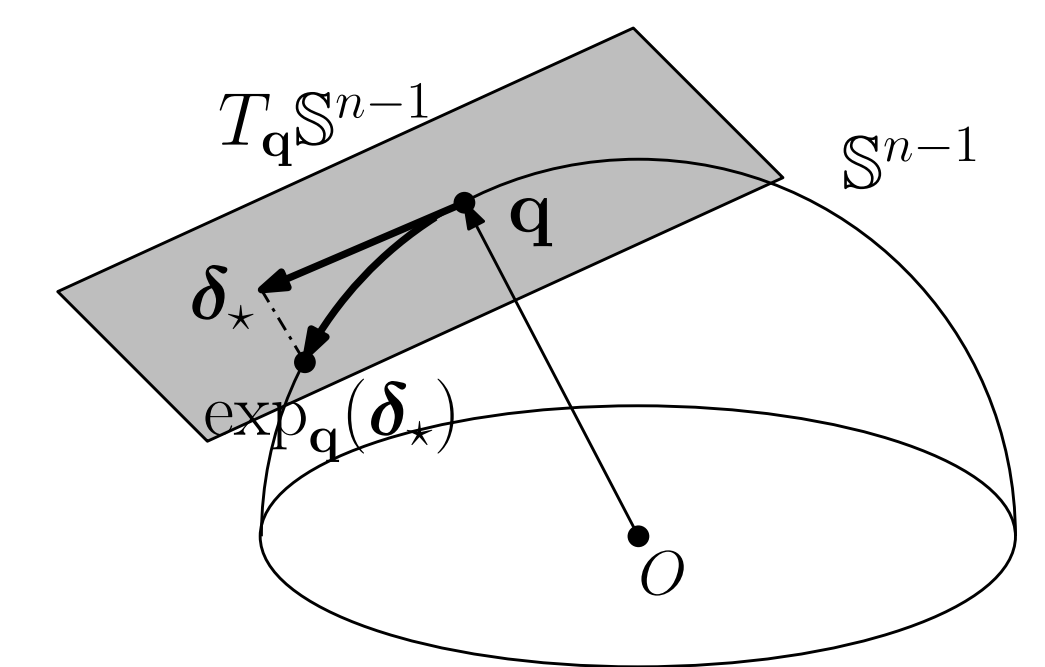
$$\|w_* - 0\| \leq \min \left\{ \frac{c_* \mu}{\theta} \sqrt{\frac{n \log p}{p}}, \frac{\mu}{16} \right\}.$$

Main Ingredients III - A Riemannian Trust-region Algorithm on Sphere

Consider $q \in \mathbb{S}^{n-1}$; for $\delta \perp q$, calculus gives

$$\begin{aligned} f(\exp_q(\delta)) &= f(q) + \langle \delta, \nabla f(q) \rangle + \frac{1}{2} \delta^* (\nabla^2 f(q) - \langle q, \nabla f(q) \rangle) \delta + O(\|\delta\|^3) \\ &= \hat{f}(\delta; q) + O(\|\delta\|^3) \end{aligned}$$

where $\exp_q(\delta) \doteq q \cos \|\delta\| + \frac{\delta}{\|\delta\|} \sin \|\delta\|$.



Basic **Riemannian trust-region method**:

$$\begin{aligned} \delta_* &\in \arg \min_{\delta \in T_{q_k} \mathbb{S}^{n-1}, \|\delta\| \leq \Delta} \hat{f}(\delta; q_k) \\ q_{k+1} &= \exp_{q_k}(\delta_*). \end{aligned}$$

The trust-region subproblem involves a (possibly nonconvex) quadratic objective and one norm constraint. **Solvable in polynomial time by root finding [More+Sorensen'83]** or SDP relaxation.

References

Ju Sun, Qing Qu, John Wright. **Complete dictionary recovery over the sphere**. Available online: <http://arxiv.org/abs/1504.06785>

